Xen is not just paravirtualization

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• Virtualization

• Xen Virtualization

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- Virtualization
- Xen Virtualization



When discussing virtualizatin ...

- 1) CPU Virtualization?
- 2) Memory Virtualization?
- 3) Device Virtualization?

• A virtual machine is taken to be an efficient, isolated duplicate of the real machine (by Formal Requirements for Virtualizable Third Generation Architectures, Gerald J.Popek and Rebert P. Goldberg, 1974)

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Trap and Emulate

- Virtual Machine (Guest) at Unprivileged Mode
- Virtual Machine Monitor (Host or Hypervisor) at Priviledged Mode



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x86 is NOT virtualizable

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 - SGDT/SIDT/SLDT, SMSW, PUSHF/POPF
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 - CALL, JMP, INT n, RET
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- Solutions:
 - Binary Translation (QEMU, VMWare)
 - Paravirtualization (Xen)
 - Hardware-Assisted Virtualization (Xen, KVM, VMWare based on Intel-VT and AMD-V)

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Solution 1/3: Binary Translation

• philosophy: rewrite critical instructions



Solution 2/3: Hardware Virtualization (Intel VT)

• philosophy: instroduce new privileged mode



- CPU hardware virtualization extensions (Intel VT or AMD-V)
- Loadable kernel module (kvm.ko, kvm-intel.ko/kvm-amd.ko)
- QEMU as userspace emulator



Solution 3/3: Paravirtualization

- philosophy: replace critical instructions with hypercalls
- A hypercall is a software trap from a domain to the hypervisor, just as a syscall is a software trap from an application to the kernel
 - x86_32: int 0x82
 - x86_64: syscall instruction
 - x86 Intel-VT vmcall instruction



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- Library Virtualization (Wine, Cygwin)



Wikipedia

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This paper presents Xen, an x86 virtual machine monitor which allows multiple commodity operating systems to share conventional hardware in a safe and resource managed fashion, but without sacrificing either performance or functionality.

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Basic Idea of Paravirtualization

Actively inform the hypervisor with the action guest is going to taken via hypercall

Xen Framework 1/2

xen hypervisor (microkernel): dictator

- scheduling, memory management, interrupt and device control
- per-domain and per-vcpu info management

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dom0 (host): privileged admin

- xm/xend/xl (libxc)
- pygrub/hvmloader
- xenstored
- qemu and paravirtual driver backend
- native device driver

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domU (guest): non-privileged user

• paravirtual driver frontend

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Convert Linux to Paravirtual Dom0/DomU

- ELF notes (Linux) or __xen_guest section (MiniOS) in kernel image
- Enable xen features in .config when building kernel

	CONFIG_XEN=Y
	CONFIG XEN DOMO=V
105 ELFNOTE(Xen, XEN_ELFNOTE_GUEST_OS, .asciz "linux")	CONETC YEN DVHVM-V
<pre>106 ELFNOTE(Xen, XEN_ELFNOTE_GUEST_VERSION, .asciz "2.6")</pre>	
107 ELFNOTE(Xen, XEN_ELFNOTE_XEN_VERSION, .asciz "xen-3.0")	CONFIG_XEN_512GB=y
108 #ifdef CONFIG_X86_32	CONFIG XEN SAVE RESTORE=y
109 ELFNOTE(Xen, XEN_ELFNOTE_VIRT_BASE, _ASM_PTRPAGE_OFFSET)	CONETG XEN BI KDEV FRONTEND=V
110 #else	CONFIC VEN BLEDEV BACKEND-m
111 ELFNOTE(Xen, XEN_ELFNOTE_VIRT_BASE,ASM_PTRSTART_KERNEL_map)	
112 /* Map the p2m table to a S12GB-aligned user address. */	CONFIG_XEN_NETDEV_FRONTEND=y
113 ELFNOTE(Xen, XEN_ELFNOTE_INIT_P2M, .quad PGDIR_SIZE)	CONFIG_XEN_NETDEV_BACKEND=m
	CONFIG INPUT XEN KBDDEV FRONTEND=m
115 ELFNOTE(Xen, XEN_ELFNOTE_ENTRY, _ASM_PIK_STATCUP_Xen)	CONFIC VEN ERREY ERONTEND m
110 ELFNOIE(Xen, XEN_ELFNOIE_HTPERCALL_PAGE, _ASM_PIK nypercall_page)	CONFIG_XEN_FBDEV_FRONTEND=M
11/ ELENDIE(AEN, AEN_ELENDIE_PEATORES, ascit iwittable_page_tables)pae_	CONFIG_XEN_BALLOON=y
110 ELEMOTE(AEII, AEN_ELEMOTE_SUPPORTED_FRATORES, CONS (FVM_FEATORES)	CONFIG XEN BALLOON MEMORY HOTPLUG=y
(1 << XENEEAT dom())	CONFIG XEN BALLOON MEMORY HOTPLUG I IMIT=512
120 (I S ACHTEAL 2000)	
122 ELENOTE(Xen, XEN_ELENOTE, LOADER, asciz "generic")	CONFIG_XEN_DEV_EVICHN=M
ELENOTE(Xen, XEN ELENOTE L1 MEN VALID.	CONFIG_XEN_BACKEND=y
124	CONFIG XEN XENBUS FRONTEND=y
125 ELFNOTE(Xen, XEN ELFNOTE SUSPEND CANCEL, .long 1)	CONFIG XEN GNTDEV=m
126 ELFNOTE(Xen, XEN ELFNOTE MOD START PFN, ,long 1)	
127 ELFNOTE(Xen, XEN_ELFNOTE_HV_START_LOW, _ASM_PTRHYPERVISOR_VIRT_START)	CONFIG_XEN_GRANI_DEV_ALLOC=M
128 ELFNOTE(Xen, XEN_ELFNOTE_PADDR_OFFSET, _ASM_PTR 0)	CONFIG_XEN_IMEM=m
129	CONFIG XEN PCIDEV BACKEND=m
130 #endif /*CONFIG_XEN */	CONETG XEN PRIVCMD=m
"arch/x86/xen/xen-head.S" 130 lines100%	

PV, HVM or PVHVM



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Xen CPU Virtualization

- vcpu \approx task_struct
- ${\, \bullet \,}$ domain \approx container or process group
- xen schedules vcpu



Event Channel Types

- Interdomain Event
- Virtual IRQ Event
- Physical IRQ Event
- IPI Event

Registration

- PVM registers event channel handler to Xen via register_callback(CALLBACKTYPE_event, xen_hypervisor_callback)
- PVHVM sets HYPERVISOR_CALLBACK_VECTOR via HYPERVISOR_hvm_op(HVMOP_set_param, &a)

Xen Interrupt Virtualization: Event Channel 2/2



Xen Hypervisor

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Xen Memory Virtualization 1/2

- Address Types
 - GVA (Guest Virtual Address)
 - GPA (Guest Physical Address) or GFN (Guest page Frame Number)
 - HPA (Host Physical Address) or MFN (Machine page Frame Number)
- Hardware-assisted Memory Virtualization (Method 1/3): Second-Level Page Table
 - : Intel: Extended Page Table (EPT)
 - : AMD: Nested Page Table (NPT)



Xen Memory Virtualization 2/2

- Direct Paging (Method 2/3): guest manage the (GVA, HPA) page table directly
- Shadow Paging (Method 3/3): xen hypervisor maintains a shadow (GVA, HPA) page table which is not awared by guest



• HVM emulated legacy device (QEMU)



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device abstraction	pci_device, pci_driver	
device discovery	PCI Tree	
device configuration	PCI Config Space (IO/MMIO)	
data flow	DMA Ring Buffer	
shared memory	N/A or IOMMU	
interrupt	IOAPIC, MSI, MSI-X	



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Xenstore/Xenbus





- Usually put grant ref (not data) in ring
- Grant ref of ring pages are shared via xenstore
- Usually one ring buffer for each device queue
- One or more pages for each ring
- Producer and Consumer (barrier)



A D > A B > A B > A

Xen Paravirtual Networking Framework



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A D > A B > A B > A B



• COLO - Coarse Grain Lock Stepping



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- LivePatch



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- Stealthy monitoring with Xen altp2m



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- Windows PV Receive Side Scaling
- More at Xen Summit and xen-devel



Publications

- Xen and the art of virtualization. Paul Barham, Boris Dragovic, Keir Fraser, Steven Hand, Tim Harris, Alex Ho, Rolf Neugebauer, Ian Pratt, and Andrew Warfield. SOSP 2003
- The Definitive Guide to the Xen Hypervisor. David Chisnall. 2007
- Intel 64 and IA-32 Architectures Software Developer Manuals
- Various system & security research paper and presentation

Miscellaneous

- Xen Project Developer Summit
- https://blog.xenproject.org
- https://github.com/finallyjustice/JOS-vmx

• What is virtualization



Image: A matrix

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- What is virtualization
- Paravirtualization and Hardware-assisted Virtualization



- What is virtualization
- Paravirtualization and Hardware-assisted Virtualization
- Xen vs. KVM



- What is virtualization
- Paravirtualization and Hardware-assisted Virtualization
- Xen vs. KVM
- Grant Table, Event Channel, Paravirtual Drivers

